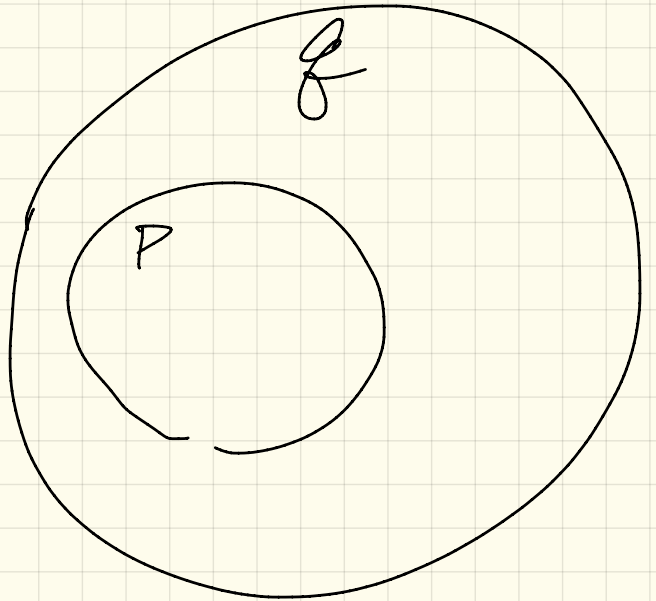
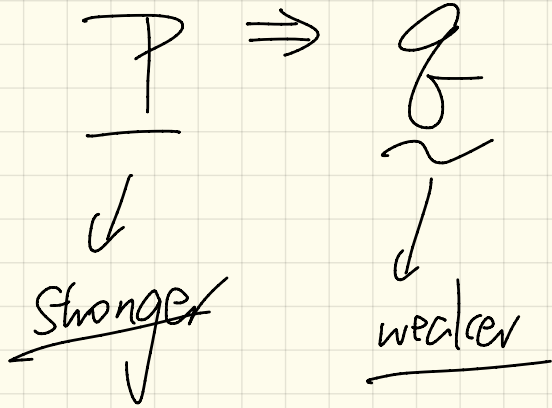


Tuesday Nov. 27

Lecture 22



# Program Correctness: Example (1)

```
class FOO
  i: INTEGER
  increment_by_9
  require
    i > 3
  do
    i := i + 9
  ensure
    i > 13
  end
end
```

$\{Q\}$      $S$      $\{R\}$

$\{i > 3\} \quad i := i + 9 \quad \{i > 13\}$

↓

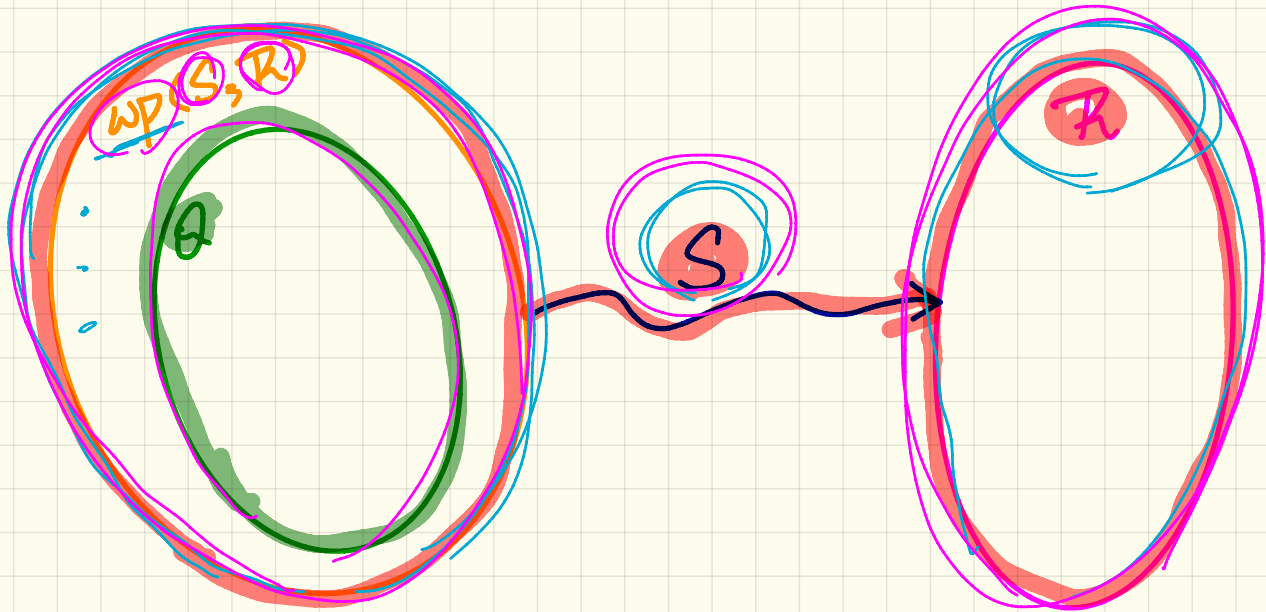
predicate  
(true or false)

## Program Correctness : Example (2)

```
class FOO
  i: INTEGER
  increment_by_9
    require
       $i > 5$ 
    do
       $i := i + 9$ 
    ensure
       $i > 13$ 
    end
end
```

# Hoare Triple as a Predicate

$$\{Q\} S \{R\} \equiv \underline{Q} \Rightarrow \underline{wp}(S, \underline{R})$$



# Program Correctness: Example (1)

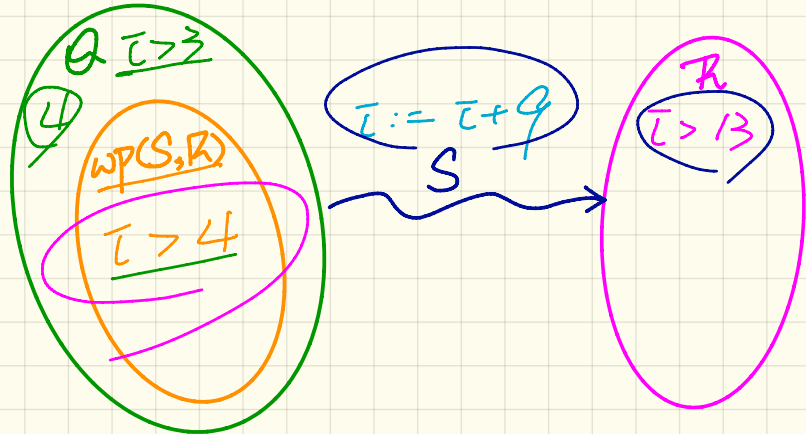
$$wp(x := x + 1, \sum_{x=1}^n x + x > 10)$$

*Annotations: x+1, free, bounded*

$$\{Q\} S \{R\} \equiv Q \Rightarrow wp(S, R)$$

```

class FOO
  i: INTEGER
  increment_by_9
  require
  4 i > 3
  do
  i := 4 i + 9
  ensure
  13 i > 13
  end
end
end
    
```



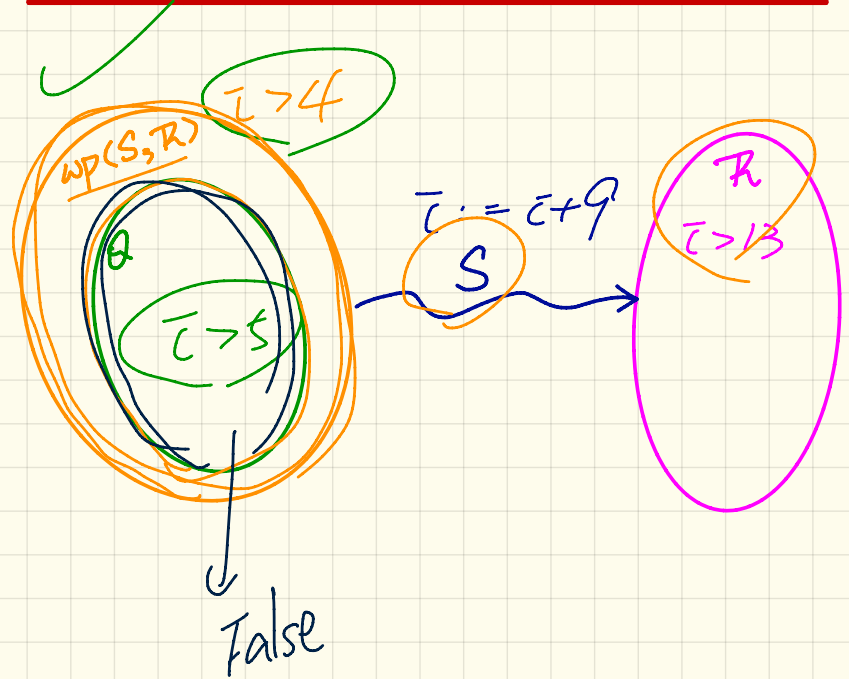
$$wp(i := i + 9, i > 13) = \{ \text{rule of wp for } := \} = i + 9 > 13$$

*Annotations: i+9, i > 13, i := i + 9*

# Program Correctness : Example (2)

```
class FOO
  i: INTEGER
  increment_by_9
  require
    i > 5
  do
    i := i + 9
  ensure
    i > 13
  end
end
```

$$\{Q\} S \{R\} \equiv Q \Rightarrow wp(S, R)$$



# wp Rules

$$wp(x := e, R) = R[x := e]$$

$$wp(\text{if } B \text{ then } S_1 \text{ else } S_2 \text{ end, } R) = \left( \begin{array}{l} B \Rightarrow wp(S_1, R) \\ \wedge \\ \neg B \Rightarrow wp(S_2, R) \end{array} \right)$$

$$wp(S_1 ; S_2, R) = wp(S_1, wp(S_2, R))$$



# Proof Rules

$$\{Q\} S \{R\} \equiv Q \Rightarrow wp(S, R)$$

$$\{Q\} x := e \{R\} \iff Q \Rightarrow \underbrace{R[x := e]}_{wp(x := e, R)}$$

$$\{Q\} \text{ if } B \text{ then } S_1 \text{ else } S_2 \text{ end } \{R\} \\ \iff \left( \begin{array}{l} \{Q \wedge B\} S_1 \{R\} \\ \wedge \\ \{Q \wedge \neg B\} S_2 \{R\} \end{array} \right) \iff \left( \begin{array}{l} (Q \wedge B) \Rightarrow wp(S_1, R) \\ \wedge \\ (Q \wedge \neg B) \Rightarrow wp(S_2, R) \end{array} \right)$$

$$\{Q\} S_1 ; S_2 \{R\} \iff Q \Rightarrow \underbrace{wp(S_1, wp(S_2, R))}_{wp(S_1 ; S_2, R)}$$

$$\text{wp} \left( \frac{x := x + 1}{x > x_0} \right)$$

$$= \{ \text{wp mke } x := \}$$

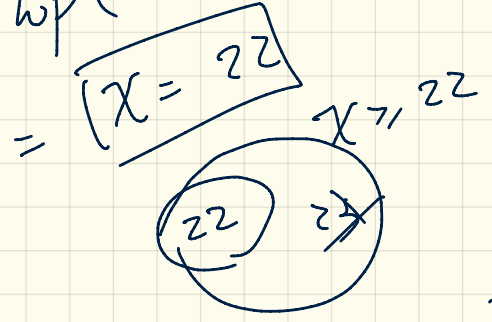
$$\underbrace{x}_{\text{post}} > \underbrace{x_0}_{\text{pre}} \left[ \underbrace{x}_{\text{post}} := \underbrace{x}_{\text{pre}} + 1 \right]$$

$$= x_0 + 1 > x_0$$



1. To be correct, any Q stronger than True.
2. To be appropriate design, up to the agreement between UC and S.

$$\text{wp} (x := x + 1, x = 23)$$



False

$wp(\text{if } B \text{ then } S_1 \text{ else } S_2 \text{ end}, R)$

$\rightarrow B \quad S_1 \quad ] \quad B \Rightarrow wp(S_1, R)$

$\neg B \quad S_2 \quad ] \quad \neg B \Rightarrow wp(S_2, R)$

# Rule of wp: Conditionals

$$wp(\text{if } B \text{ then } S_1 \text{ else } S_2 \text{ end}, R)$$

$$B \Rightarrow wp(S_1, R)$$

$$\vee$$

$$\neg B \Rightarrow wp(S_2, R)$$

$$B \Rightarrow wp(S_1, R)$$

$$\wedge$$

$$\neg B \Rightarrow wp(S_2, R)$$

??

vs.  
choose  
 $x = -1$   
 $x = -1$

Consider:  $x \geq -1$   
 $\vee$   
 $x \geq 1$

$$x \geq -1$$

$$x \geq -1$$

$$\wedge$$

$$x \geq 1$$

$$\equiv x \geq 1$$

$$wp(\text{if } y > 0 \text{ then } x := x + 1 \text{ else } x := x - 1 \text{ end}, x \geq 0)$$

$$y > 0 \Rightarrow wp(x := x + 1, x \geq 0) \quad x \geq -1$$

$$y \leq 0 \Rightarrow wp(x := x - 1, x \geq 0) \quad x \geq 1$$

{R}

if B<sub>1</sub> then

S<sub>1</sub>

elseif B<sub>2</sub> then

S<sub>2</sub>

else

S<sub>3</sub> B<sub>1</sub>  $\Rightarrow$  wp (S<sub>3</sub>, R)

end

{R}

$\neg B_1 \wedge B_2 \Rightarrow$  wp (S<sub>2</sub>, R)

$\neg (B_1 \vee B_2) \Rightarrow$  wp (S<sub>3</sub>, R)

{Q} if B then S1 else S2 end {R}

$$\underline{Q} \Rightarrow \left( \begin{array}{l} \neg B \Rightarrow \text{wp}(S_1, R) \\ \wedge \\ \neg B \Rightarrow \text{wp}(S_2, R) \end{array} \right)$$

shunting

$$\begin{array}{l} P \Rightarrow (Q \Rightarrow R) \\ \equiv \\ (P \wedge Q) \Rightarrow R \end{array}$$

# Correctness of Program: Conditionals

Is this program correct?

```
{x > 0 ∧ y > 0}
if x > y then
  bigger := x ; smaller := y
else
  bigger := y ; smaller := x
end
{bigger ≥ smaller}
```

1. Calculate the wp

wp (if  $x > y$  then  $b := x ; s := y$  else  $b := y ; s := x$  end,  $b \geq s$ )

2. Prove or disprove

$x > 0 \wedge y > 0 \Rightarrow \text{wp}$

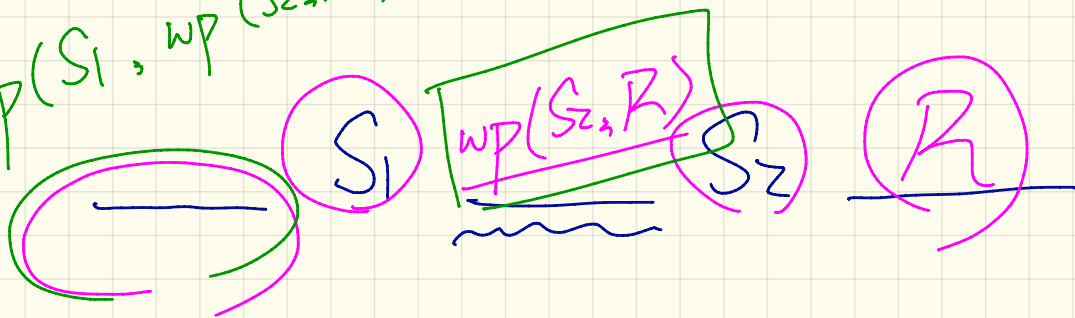
$= x > y \Rightarrow \text{wp}(\text{[ ]}, b \geq s)$

$\neg(x > y) \Rightarrow \text{wp}(\text{[ ]}, b \geq s)$

$$\text{wp}(S_1) \equiv \boxed{S_2, R}$$

$$= \text{wp}(S_1, \underline{\text{wp}(S_2, R)})$$

$$\text{wp}(S_1, \text{wp}(S_2, R))$$





# Correctness of Program: Sequential Composition

Is  $\{ \text{True} \} \text{tmp} := x; x := y; y := \text{tmp} \{ x > y \}$  correct?

Q: Swap  $x$  and  $y$

without using a temp variable.

$$\begin{aligned} x &= x + y - x \\ y &= \end{aligned}$$

$$\begin{aligned} x &= x + y \\ y &= y - x \\ x &= y + x \\ y &= y \end{aligned}$$